

1 From Big Computer Algorithms to Small Computer Algorithm

1.1 Strategy

Size Reduction + Recursion + Terminate when brute force size is sufficiently small

1.2 Simple idea

Choose a split, divide the set into two part and cost time to work on the part contains median

2 Find k-th Smallest Element in Big Computer

Question:

Ranking sequence $A(n, k)$ ($1 \leq k \leq n$), find k-th smallest element in A of n elements

2.1 Algorithms

1. **randomly** pick a element of A;
2. use s as the splitter to divide A into $B = A_{smaller}(s)$, $C = A_{bigger}(s)$
3. suppose $b = |B| \leq k$
 find $C(n-b, k-b)$
 else $b > k$,
 find $B(b, k)$

Emphasizing: RANDAMLY!

2.2 Time

consider expected time:

$$T(n) = \left(\sum_{b=1}^{k-1} T(n-b) + \sum_{b=k+1}^{n-1} T(b) \right) \times \frac{1}{n-1} + n$$

solution:

$$T(n) = O(n)$$

2.3 Proof

Decide what is d for the proof to hold:

Assume

$$T(m) \leq d \times m$$

For all $T(m) \leq d \times m$

$$T(n) = \left(\sum_{b=1}^{k-1} T(n-b) + \sum_{b=k+1}^{n-1} T(b) \right) \times \frac{1}{n-1} \leq d \times \left(\sum_{b=1}^{k-1} T(b) \right) \times \frac{1}{n-1} + 1 \leq d \times n$$

3 Find Median in Small Computer

Questions:

How to find median of A ? (no enough space for every elements, one pass algorithm, which means every element goes into cache once)

Let the size of cache be F , how can we find the median?

3.1 Method

Throw away max or min or both with a high probability

- If the new number is larger than $\max(S)$ or smaller than $\min(S)$ remove it to place in H or L accordingly
- If the new number is in $(\min(S), \max(S))$, then keep it and remove $\max(S)$ or $\min(S)$ to make L or H more balanced

3.2 Algorithms

1. take $|F|$ elements median in randomly (every element in is chosen at **random**)
2. throw away an item (min or max) (median, though not know which is kept with high probability compete it)
3. redo step 2, until every item in F , L or H
4. sort (F), return the median

BEST situation when the size of F is \sqrt{n}